## 1 A construction with straightedge alone

A problem I spotted on someone's Math 241 homework (paraphrased):
Given a circle, its diameter $A B$, and a point $P$ lying neither on the circle nor on the given diameter extended, construct with straightedge alone a line passing through $P$ and perpendicular to $A B$ (extended).
(The traditional solution using both compass and straightedge is Euclid I.12.)
Since the only thing a straightedge can do is join points to form lines (and thereby find new points as intersections), there's only so much you can try; so it's not hard to find the right construction.


Construction Join $P A$, intersecting the circle also at $A^{\prime}$. Join $P B$, intersecting the circle also at $B^{\prime}$. Join $A^{\prime} B$ and $A B^{\prime}$, and extend them to intersect at $Q$. Join $P Q$. Then $P Q \perp A B$, as desired.

Proof $A B$ is a diameter, so $\angle A A^{\prime} B$ and $\angle A B^{\prime} B$ are right angles (Euclid III.31). Thus the segments $\mathrm{Q}^{\prime}$ and $\mathrm{PB}^{\prime}$ are altitudes in $\triangle A P Q$, and so their point of intersection, $B$, is the orthocentre. Therefore $A B$ extended is the third altitude, and as such forms a right angle with PQ , as claimed.
(This construction ignores some special cases. Identifying them and patching things up is left as an exercise.)

## 2 Why sine isn't a polynomial

A few reasons why sine isn't a polynomial: it has infinitely many zeroes, but isn't the zero polynomial; its Maclaurin series isn't finite; it has infinitely many nonzero derivatives; it is bounded but not constant. Finally, a nifty one, exercise 1.1.13 from Polynomials, by E.J. Barbeau:

Find all polynomials $f$ such that $f(2 t)$ can be written as a polynomial in $f(t)$, i.e., for which there exists a polynomial $h$ such that

$$
f(2 t)=h(f(t))
$$

Use the identity $\sin ^{2} 2 t=4 \sin ^{2} t\left(1-\sin ^{2} t\right)$ to show that $\sin t$ is not a polynomial.

This one will go on the list of outstanding problems.

## 3 Making or

Last time we discussed our outstanding problem of showing that it's impossible to express $\vee$ using only $\neg$ and $\underline{\vee}$. The notes for 2006 Nov 20 include a solution for the similar problem of showing it's impossible to express $\neg$ using only $\wedge, \vee$, and $\underline{\vee}$; we defined the set of boolean functions expressible with these operators recursively, then showed by structural induction that they all "fix zero", and observed that negation doesn't.

To find a similar proof for the problem of expressing $\vee$, we listed all sixteen boolean functions of two boolean variables $A$ and $B$, and tried to express them all using $\neg$ and $\underline{\vee}$.

| 0 | 0 | 1 | 1 | $A$ |
| :--- | :--- | :--- | :--- | :--- |
| 0 | 1 | 0 | 1 | $B$ |
| 0 | 0 | 0 | 0 | $A \underline{\vee} A$ |
| 0 | 0 | 0 | 1 |  |
| 0 | 0 | 1 | 0 |  |
| 0 | 0 | 1 | 1 | $A$ |
| 0 | 1 | 0 | 0 |  |
| 0 | 1 | 0 | 1 | $B$ |
| 0 | 1 | 1 | 0 | $A \underline{\vee} B$ |
| 0 | 1 | 1 | 1 |  |


| 0 | 0 | 1 | 1 | $A$ |
| :--- | :--- | :--- | :--- | :--- |
| 0 | 1 | 0 | 1 | $B$ |
| 1 | 0 | 0 | 0 |  |
| 1 | 0 | 0 | 1 | $\neg(A \underline{B})$ |
| 1 | 0 | 1 | 0 | $\neg B$ |
| 1 | 0 | 1 | 1 |  |
| 1 | 1 | 0 | 0 | $\neg A$ |
| 1 | 1 | 0 | 1 |  |
| 1 | 1 | 1 | 0 |  |
| 1 | 1 | 1 | 1 | $\neg(A \underline{\vee})$ |

It's pretty easy to find the above eight expressions. After fiddling for a while, it's easy to become convinced that they're the only ones that can be expressed using just $\neg$ and $\underline{\vee}$. (Other candidate expressions, such as $(\neg A) \underline{\vee} B$, turn out to be equivalent to ones already listed.)

After staring at the table for a while, we notice what distinguishes those eight expressions from the other eight - the expressible functions have an even number of 1 s in their rows.

The simplest way to express this condition formally is:

$$
\begin{equation*}
f(0,0) \underline{\vee} f(0,1) \underline{\vee}(1,0) \underline{\vee} f(1,1)=0 \tag{*}
\end{equation*}
$$

(Note that $\underline{V}$ is associative, so we don't need parentheses here.) So, to show that $\vee$ cannot be expressed using only $\neg$ and $\underline{\vee}$, we must show four things:

1. The functions $(A, B) \mapsto A$ and $(A, B) \mapsto B$ have property $(*)$.
2. If $f$ has property $(*)$, then $\neg f$ has it too.
3. If $f$ and $g$ have property $(*)$, then $f \underline{\vee} g$ has it too.
4. The function $(A, B) \mapsto A \vee B$ does not have property $(*)$.

I leave the details as an exercise.

