Math Club Notes: 2006 November 27

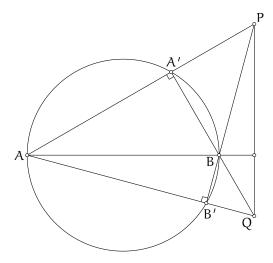
1 A construction with straightedge alone

A problem I spotted on someone's Math 241 homework (paraphrased):

Given a circle, its diameter AB, and a point P lying neither on the circle nor on the given diameter extended, construct with straightedge alone a line passing through P and perpendicular to AB (extended).

(The traditional solution using both compass and straightedge is Euclid I.12.)

Since the only thing a straightedge can do is join points to form lines (and thereby find new points as intersections), there's only so much you can try; so it's not hard to find the right construction.



Construction Join PA, intersecting the circle also at A'. Join PB, intersecting the circle also at B'. Join A'B and AB', and extend them to intersect at Q. Join PQ. Then PQ \perp AB, as desired.

Proof AB is a diameter, so $\angle AA'B$ and $\angle AB'B$ are right angles (Euclid III.31). Thus the segments QA' and PB' are altitudes in $\triangle APQ$, and so their point of intersection, B, is the orthocentre. Therefore AB extended is the third altitude, and as such forms a right angle with PQ, as claimed.

(This construction ignores some special cases. Identifying them and patching things up is left as an exercise.)

2 Why sine isn't a polynomial

A few reasons why sine isn't a polynomial: it has infinitely many zeroes, but isn't the zero polynomial; its Maclaurin series isn't finite; it has infinitely many nonzero derivatives; it is bounded but not constant. Finally, a nifty one, exercise 1.1.13 from *Polynomials*, by E.J. Barbeau:

Find all polynomials f such that f(2t) can be written as a polynomial in f(t), i.e., for which there exists a polynomial h such that

$$f(2t) = h(f(t))$$

Use the identity $\sin^2 2t = 4 \sin^2 t (1 - \sin^2 t)$ to show that sin t is not a polynomial.

This one will go on the list of outstanding problems.

3 Making or

Last time we discussed our outstanding problem of showing that it's impossible to express \lor using only \neg and \trianglelefteq . The notes for 2006 Nov 20 include a solution for the similar problem of showing it's impossible to express \neg using only \land , \lor , and \trianglelefteq ; we defined the set of boolean functions expressible with these operators recursively, then showed by structural induction that they all "fix zero", and observed that negation doesn't.

To find a similar proof for the problem of expressing \lor , we listed all sixteen boolean functions of two boolean variables A and B, and tried to express them all using \neg and \lor .

0	0	1	1	A	0	0	1	1	A
0	1	0	1	В	0	1	0	1	В
0	0	0	0	A⊻A	1	0	0	0	
0	0	0	1		1	0	0	1	$\neg(A \lor B)$
0	0	1	0		1	0	1	0	$\neg B$
0	0	1	1	A	1	0	1	1	
0	1	0	0		1	1	0	0	$\neg A$
0	1	0	1	В	1	1	0	1	
0	1	1	0	A⊻B	1	1	1	0	
0	1	1	1		1	1	1	1	$\neg(A \lor A)$

It's pretty easy to find the above eight expressions. After fiddling for a while, it's easy to become convinced that they're the only ones that can be expressed using just \neg and \checkmark . (Other candidate expressions, such as $(\neg A) \checkmark B$, turn out to be equivalent to ones already listed.)

After staring at the table for a while, we notice what distinguishes those eight expressions from the other eight — the expressible functions have an even number of 1s in their rows.

The simplest way to express this condition formally is:

$$f(0,0) \leq f(0,1) \leq f(1,0) \leq f(1,1) = 0.$$
 (*)

(Note that \forall is associative, so we don't need parentheses here.) So, to show that \lor cannot be expressed using only \neg and \forall , we must show four things:

- 1. The functions $(A, B) \mapsto A$ and $(A, B) \mapsto B$ have property (*).
- 2. If f has property (*), then \neg f has it too.
- 3. If f and g have property (*), then $f \lor g$ has it too.
- 4. The function $(A, B) \mapsto A \lor B$ does not have property (*).

I leave the details as an exercise.

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